



Hamilton Institute



**Maynooth
University**
National University
of Ireland Maynooth

It's not about DNA, It's all about Pizza

Ahmed Shalaby

2nd year PhD

Supervisor: Damien Woods

CS workshop [19/06/2024]



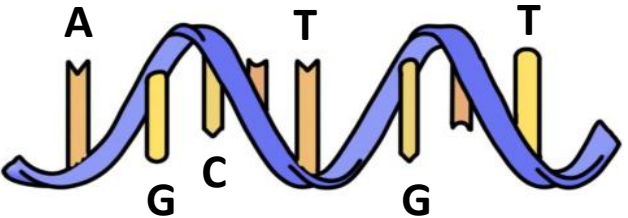
European
Innovation
Council



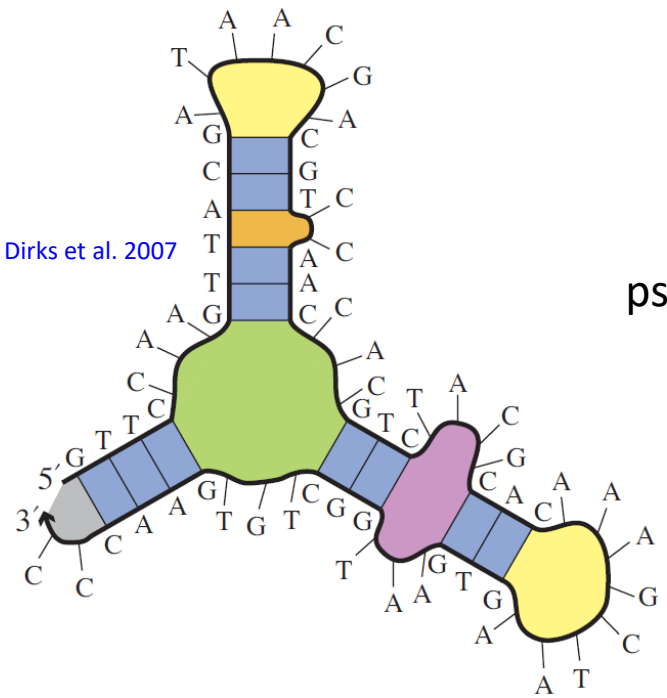
Funded by
the European Union



Secondary structure



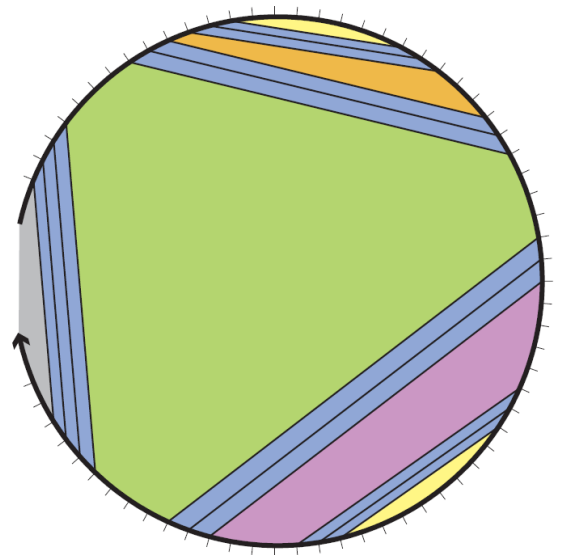
Single stranded DNA



Dirks et al. 2007

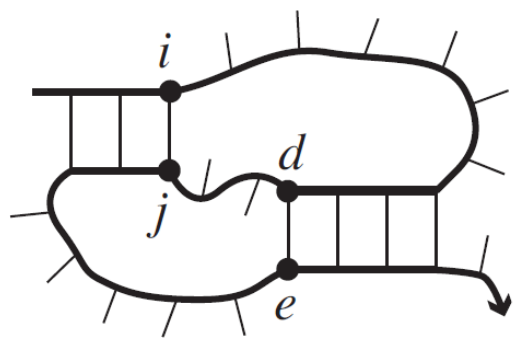
Secondary structure

pseudoknot-free



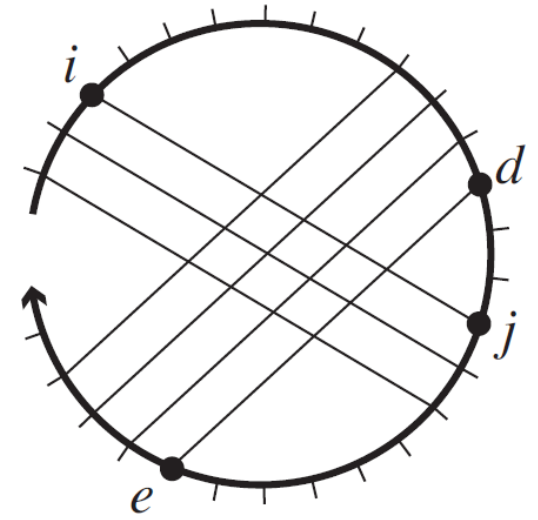
Dirks et al. 2007

Polymer graph representation



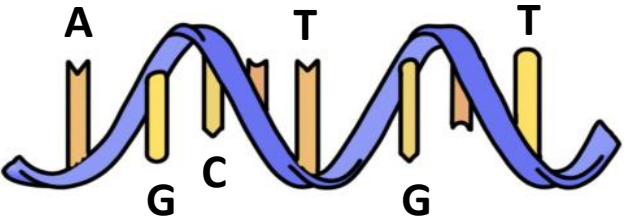
Dirks et al. 2007

pseudoknotted

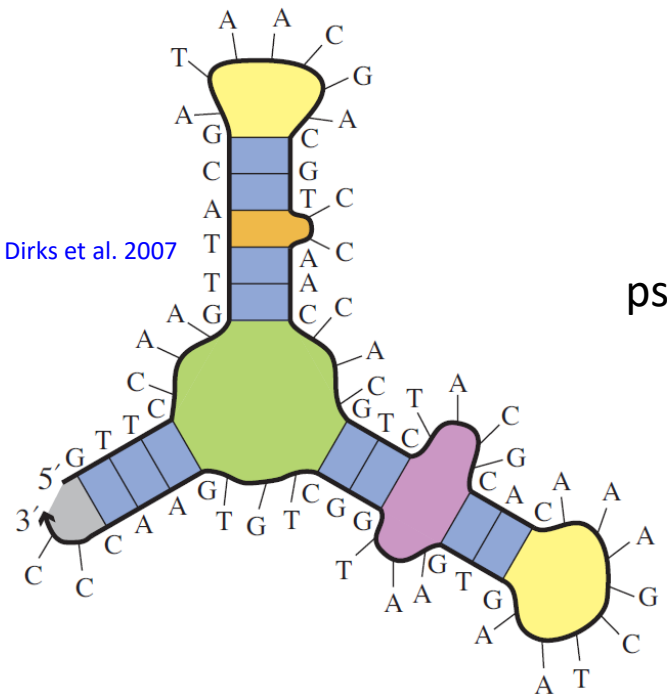


Dirks et al. 2007

Secondary structure



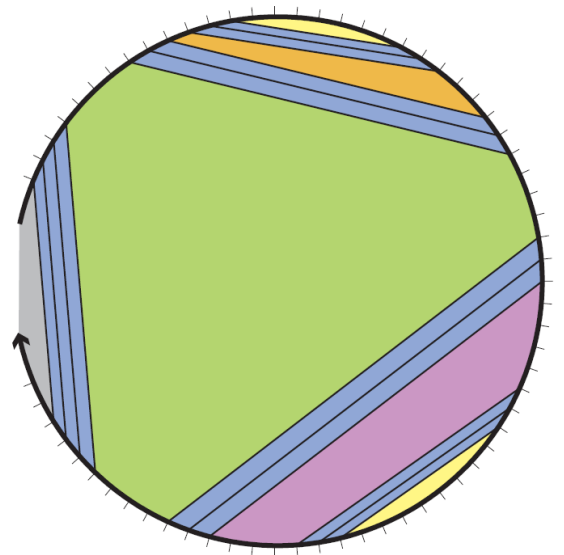
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Secondary structure

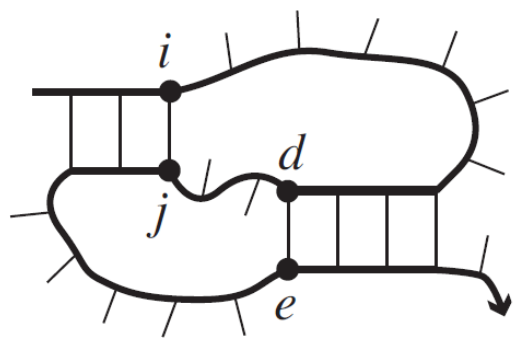
pseudoknot-free



Dirks et al. 2007

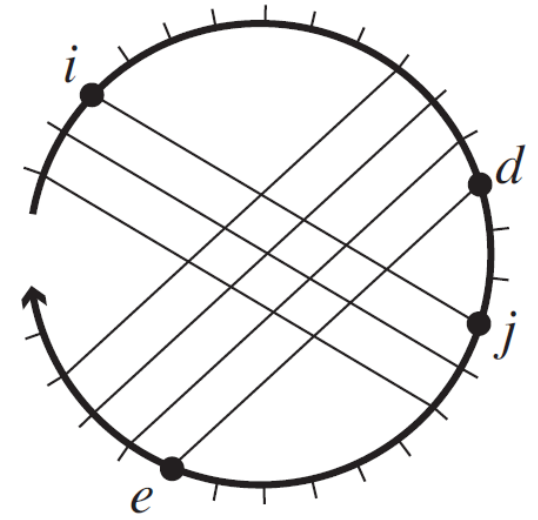
Polymer graph representation

NP – Complete



Dirks et al. 2007

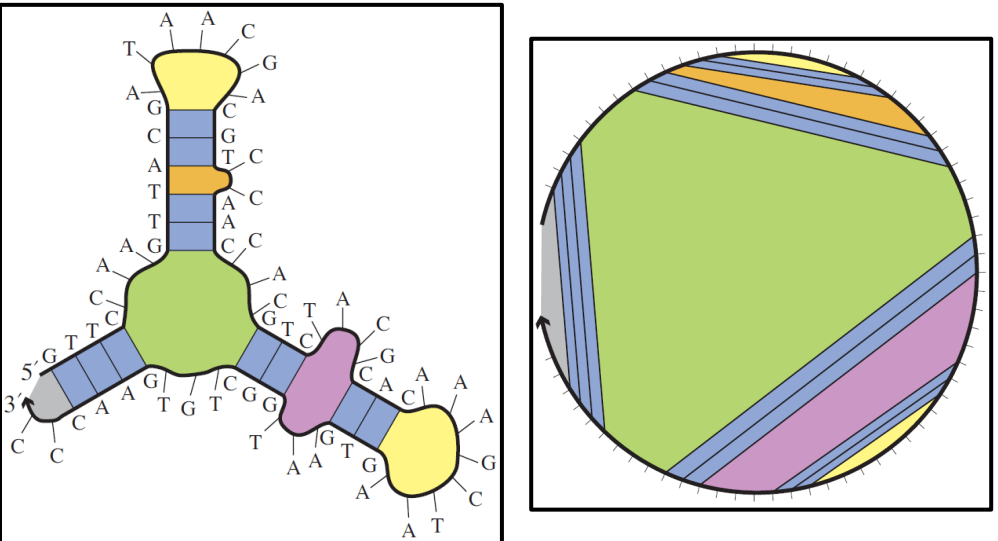
pseudoknotted



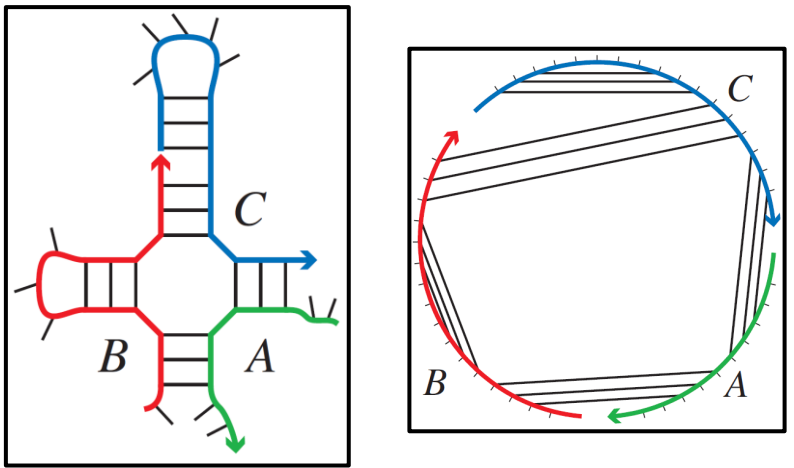
Dirks et al. 2007

Energy models, Minimum Free Energy and Partition Function

Single stranded system

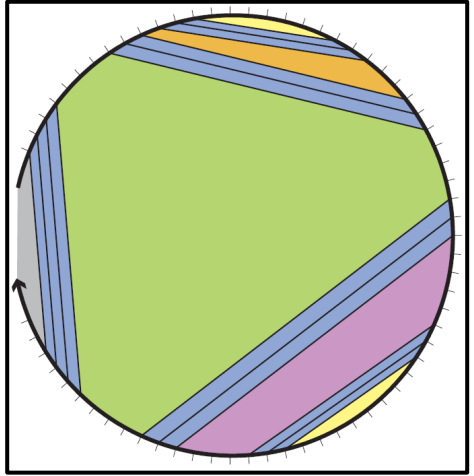
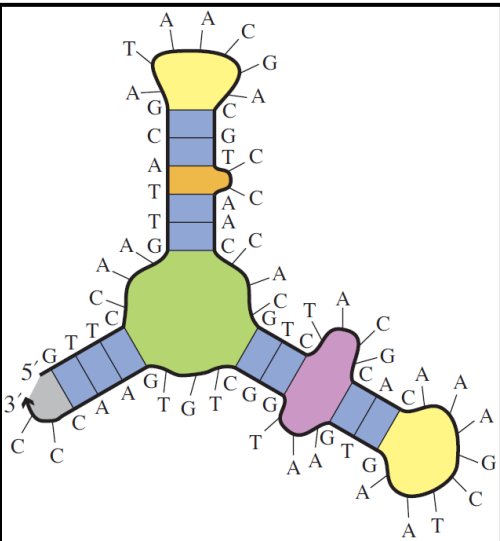


Multi stranded system of s strands



Energy models, Minimum Free Energy and Partition Function

Single stranded system

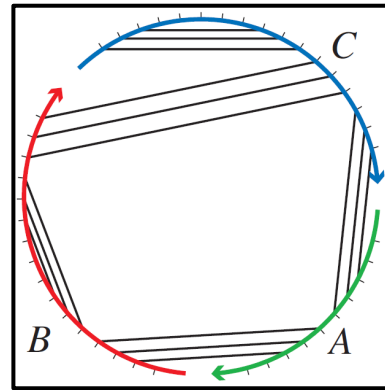
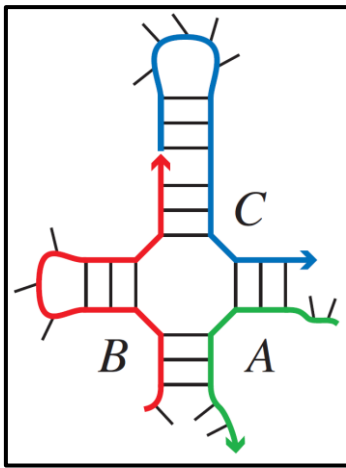


$$\Delta G(S)$$

Energy model

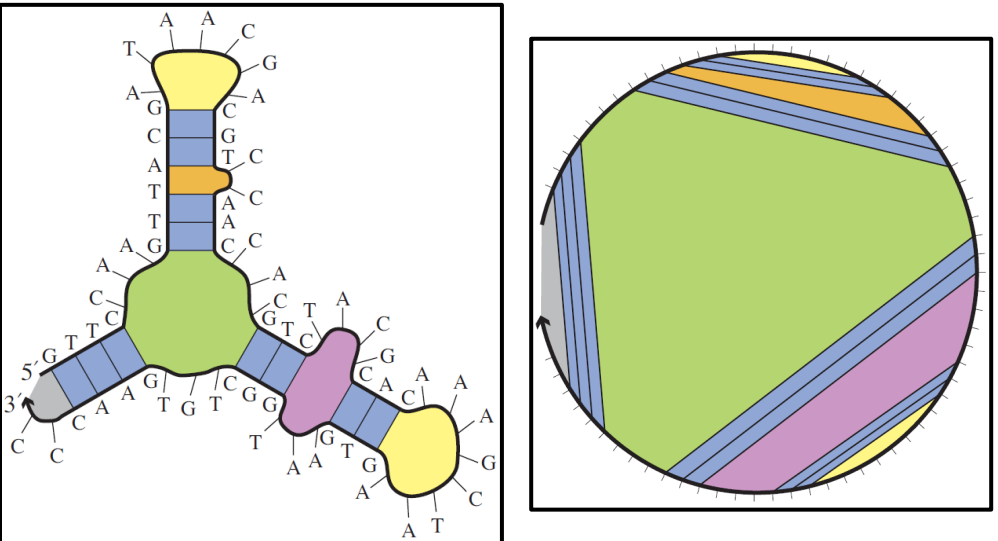
Capture the free energy of secondary structure

Multi stranded system of s strands

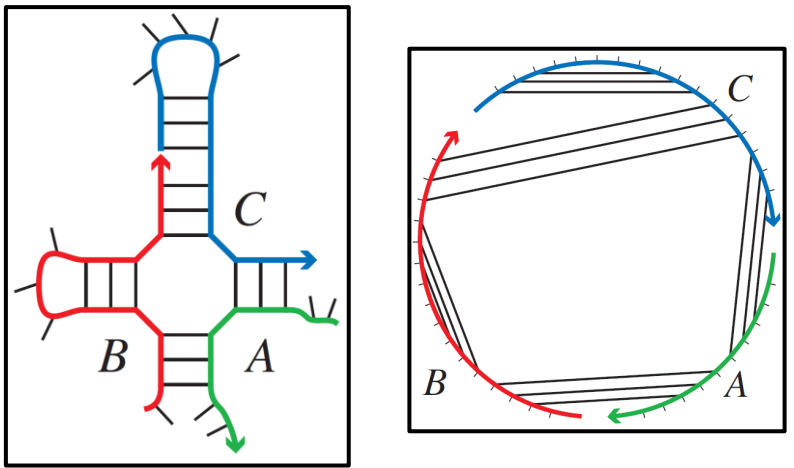


Energy models, Minimum Free Energy and Partition Function

Single stranded system



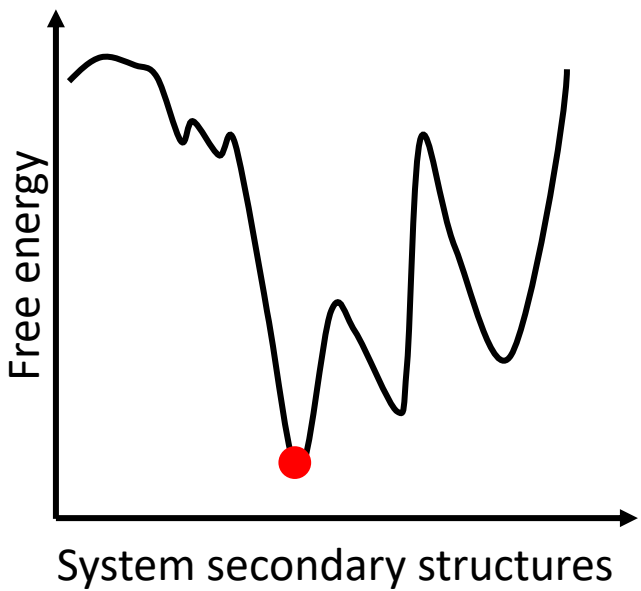
Multi stranded system of s strands



$$\Delta G(S)$$

Energy model

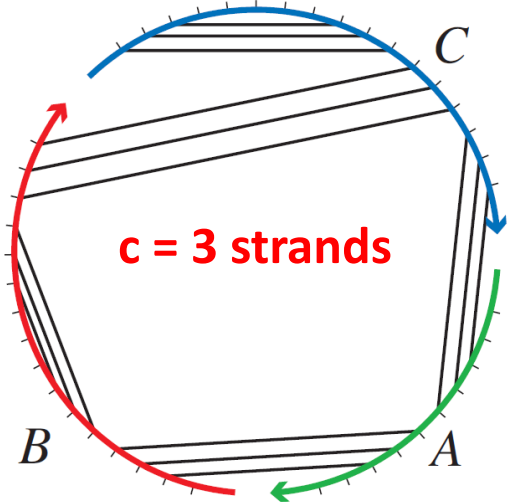
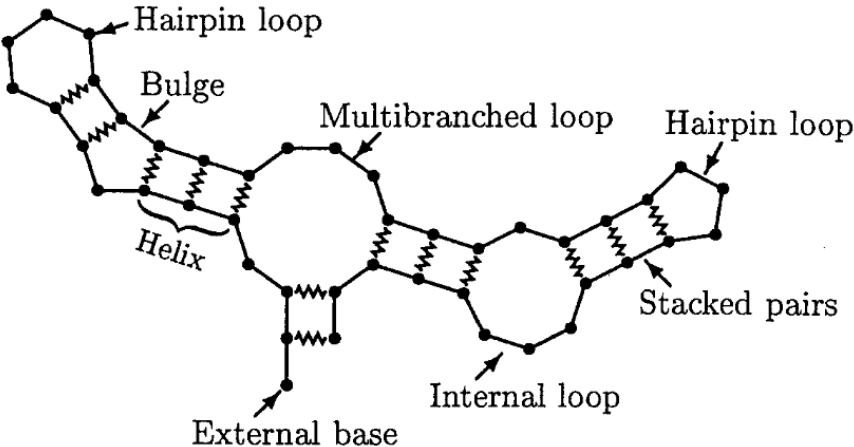
Capture the free energy of secondary structure



$$\text{MFE} = \min_{S \in \Omega} \Delta G(S)$$

Minimum Free Energy

Energy model: Loop model



$$\Delta G(S) = \sum_{l \in S} \Delta G(l)$$

$$\Delta G(S) = \sum_l \Delta G(l) + (c - 1) \Delta G^{\text{assoc}}$$

$$\min_{S \in \Omega} \Delta G(S)$$

Ω : the set of all secondary structures

Energy model: Loop model (allowing repeats)

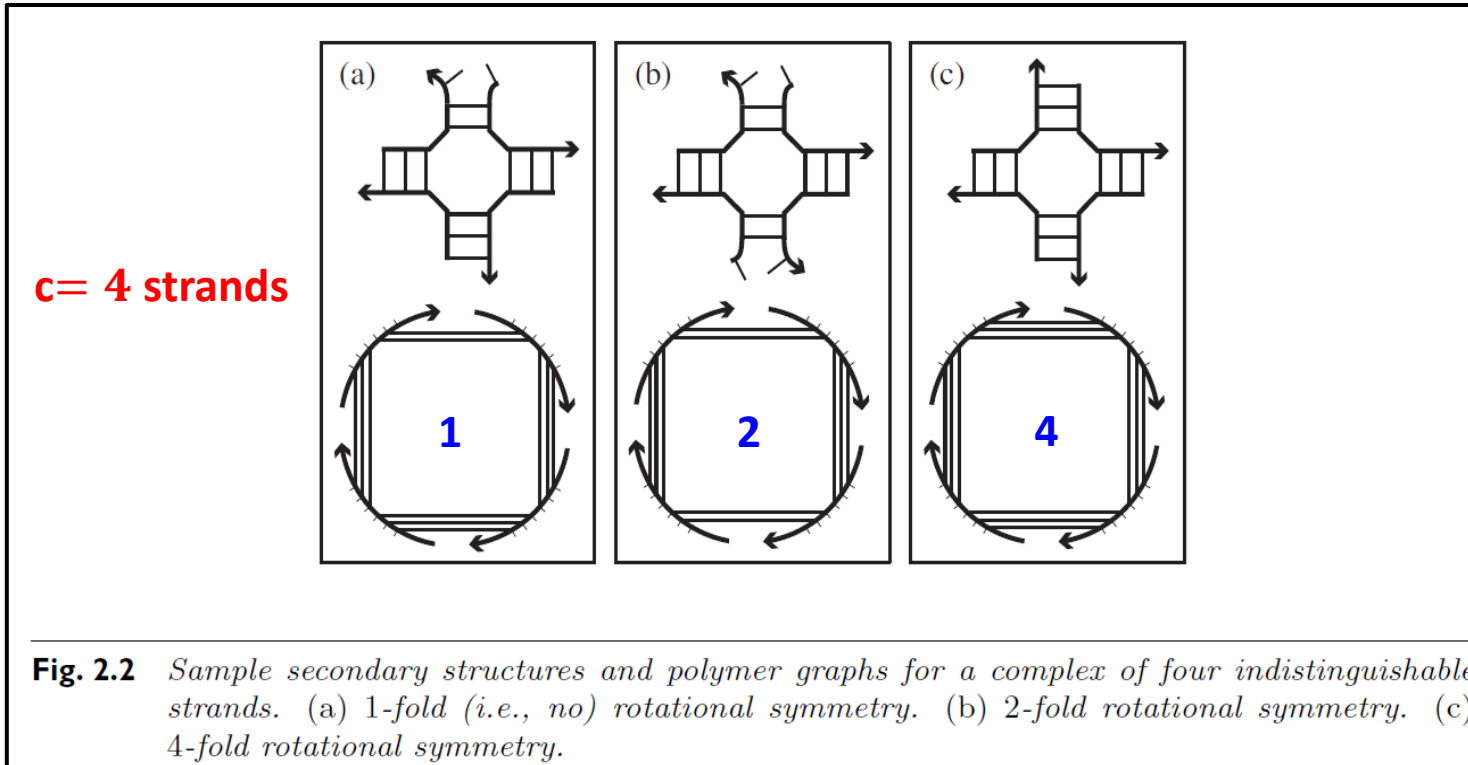
Free energy

Loop energy

Entropic association cost

Symmetry penalty

$$\Delta G(S) = \sum_{l \in S} \Delta G(l) + (c - 1) * \Delta G^{\text{assoc}} + k_B T * \log R$$



$$\min_{S \in \Omega} \Delta G(S)$$

Ω : the set of all connected structures
 R : degree of rotational symmetry

Computational complexity of Minimum Free Energy algorithms

Input Type	MFE
Single Strand	$O(N^3)$
Multiple unique Strands, Bounded ($\leq c$)	$O(N^3(c-1)!)$
Multiple Strands allowing repeats , Bounded ($\leq c$)	?
Multiple Strands, Unbounded	NP – Complete

N bases, c strands

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N bases, c strands

Open problem for ≈ 20 years

Why symmetry makes it difficult?

Input Type	MFE
Single Strand (Loop model)	$O(N^3)$
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} **Dynamic programming algorithms**

N bases, c strands

All of these are **dynamic programming** algorithms

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} **Dynamic programming algorithms**

N bases, c strands

All of these are **dynamic programming** algorithms

Subproblems  Big problem

Level	Input Type	MFE
1	Single Strand (Maximum matching)	$O(N^3)$
2	Single Strand (Loop model)	$O(N^3)$
3	Multiple unique Strands, Bounded ($\leq c$)	$O(N^3(c-1)!)$
4	Multiple Strands, Bounded ($\leq c$)	?

*N bases, c strands

All of these are **dynamic programming** algorithms

Subproblems



Big problem

Level	Input Type	MFE
1	Single Strand (Maximum matching)	$O(N^2)$
2	Single Strand (Loop model)	$O(N^3)$
3	Multiple unique Strands, Bounded ($\leq c$)	$O(N^3(c-1)!)$
4	Multiple Strands, Bounded ($\leq c$)	?

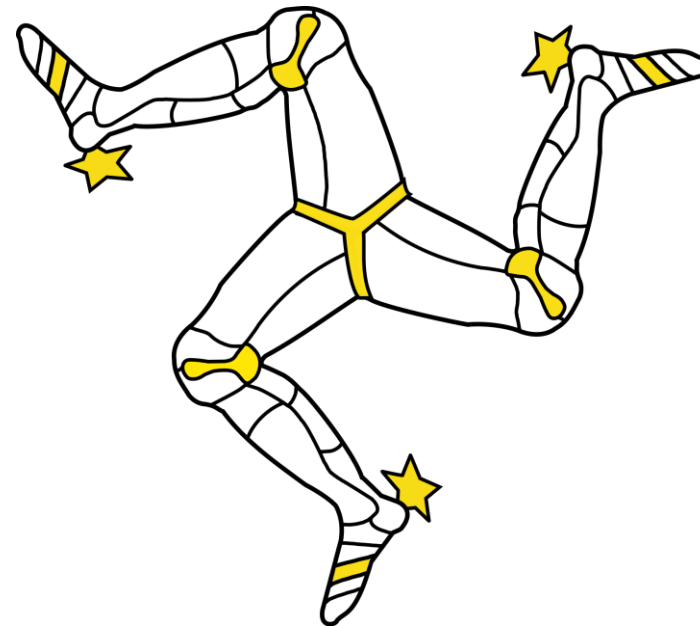
*N bases, c strands

All of these are **dynamic programming algorithms**

Subproblems



Big problem

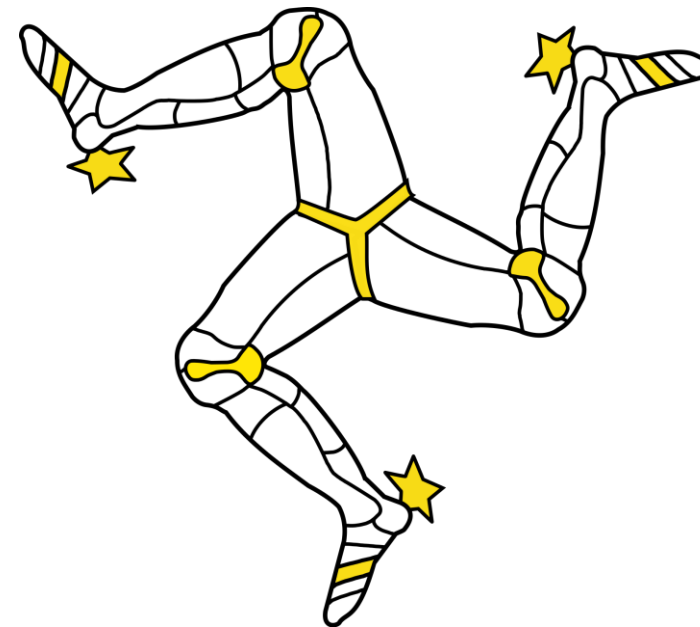


Level	Input Type	MFE
1	Single Strand (Maximum matching)	$O(N^3)$
2	Single Strand (Loop model)	$O(N^3)$
3	Multiple unique Strands, Bounded ($\leq c$)	$O(N^3(c-1)!)$
4	Multiple Strands, Bounded ($\leq c$)	?

*N bases, c strands

All of these are **dynamic programming algorithms**

Subproblems \longrightarrow Big problem



Global property

Level	Input Type	MFE
1	Single Strand (Maximum matching)	$O(N^3)$
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4	Multiple Strands, Bounded ($\leq c$)	?

*N bases, c strands

All of these are **dynamic programming algorithms**

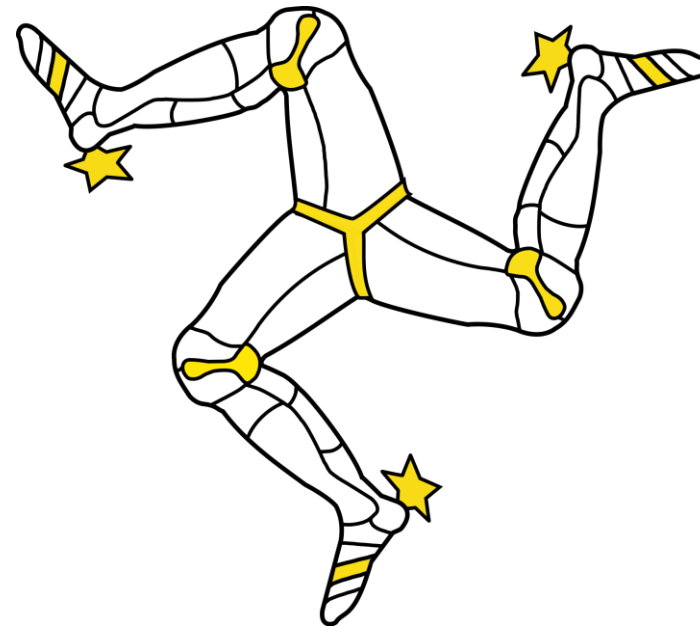
Subproblems



Big problem



Local point of view



Global property

Level	Input Type	MFE
1	Single Strand (Maximum matching)	$O(N^3)$
2	Single Strand (Loop model)	$O(N^3)$
3	Multiple unique Strands, Bounded ($\leq c$)	$O(N^3(c-1)!)^c$
4	Multiple Strands, Bounded ($\leq c$)	?

³ N bases, c strands

All of these are **dynamic programming algorithms**

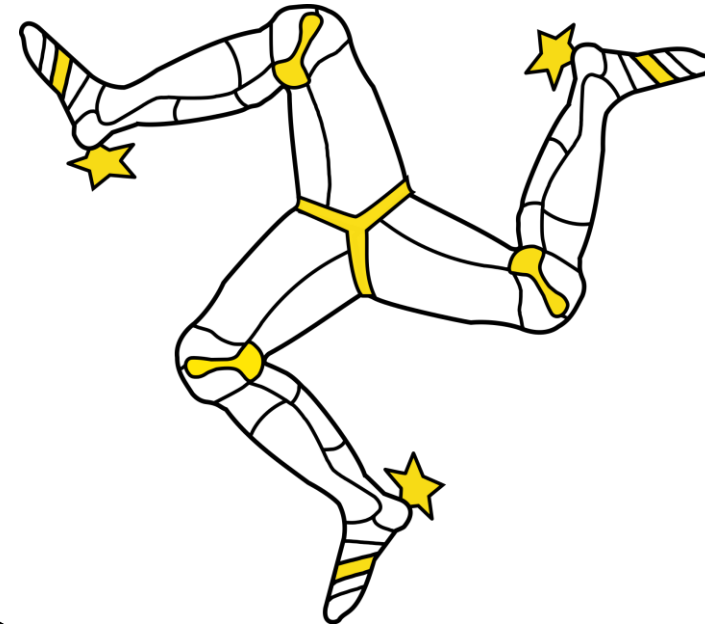
Subproblems



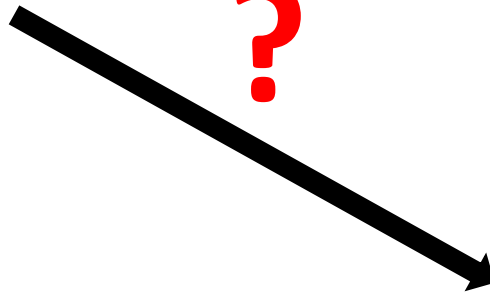
Big problem



Local point of view



?



Global property

Possible approach

Input Type	MFE
Single Strand (Loop model)	$O(N^3)$
Multiple unique Strands, Bounded ($\leq c$)	$O(N^3(c - 1)!)$
Multiple Strands, Bounded ($\leq c$)	?

N bases, c strands

$$\Delta G(S) = \sum_l \Delta G(l) + (c - 1)\Delta G^{\text{assoc}} + k_B T * \log R$$

Possible approach

Input Type	MFE
Single Strand (Loop model)	$O(N^3)$
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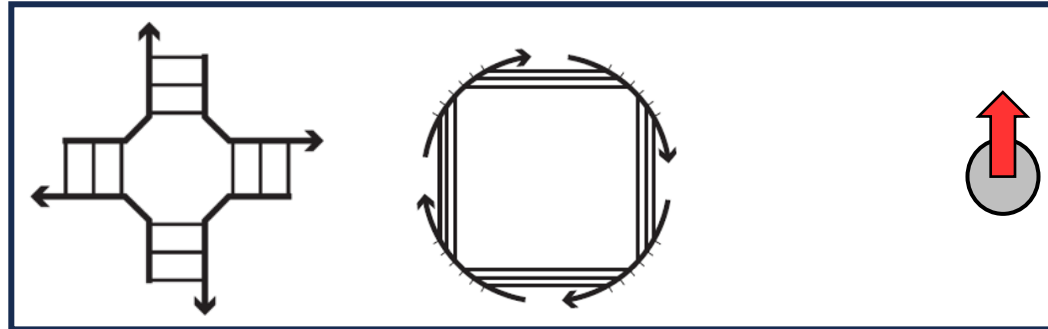
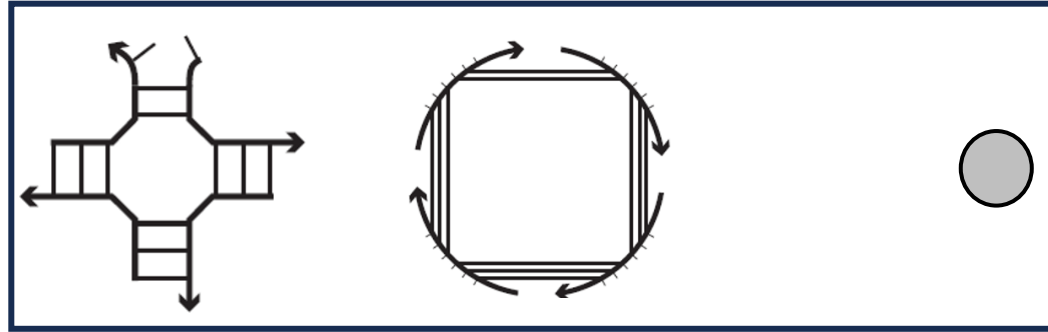
$$\Delta G(S) = \sum_l \Delta G(l) + (c - 1)\Delta G^{\text{assoc}} + k_B T * \log R$$

$$\Delta G'(S) = \sum_l \Delta G(l) + (c - 1) \Delta G^{\text{assoc}}$$

Ignore symmetry

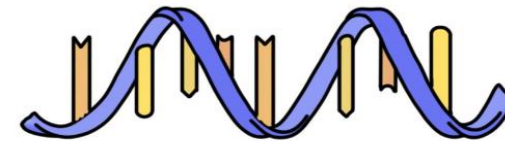
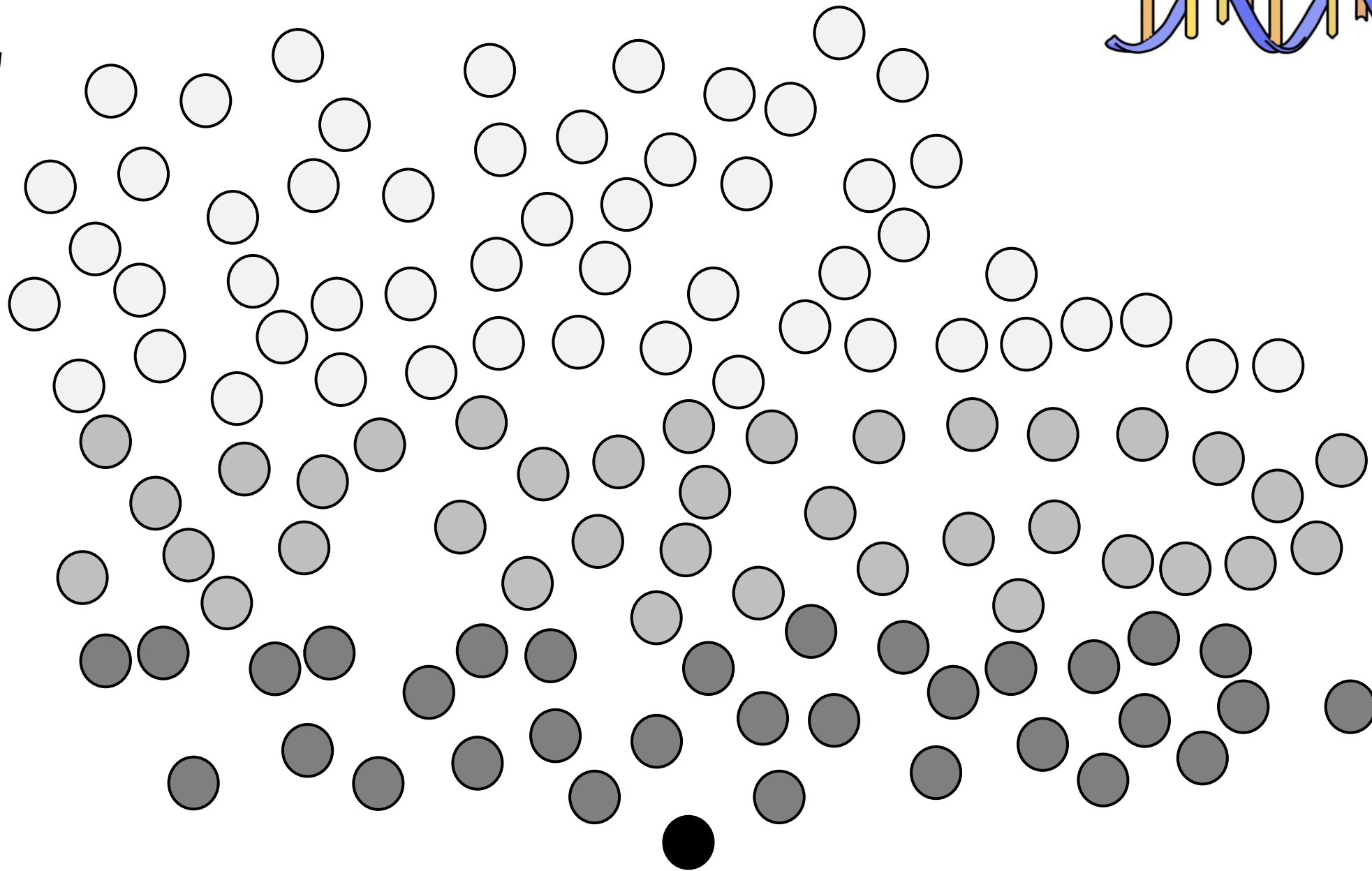
Let's ignore the symmetry for a while

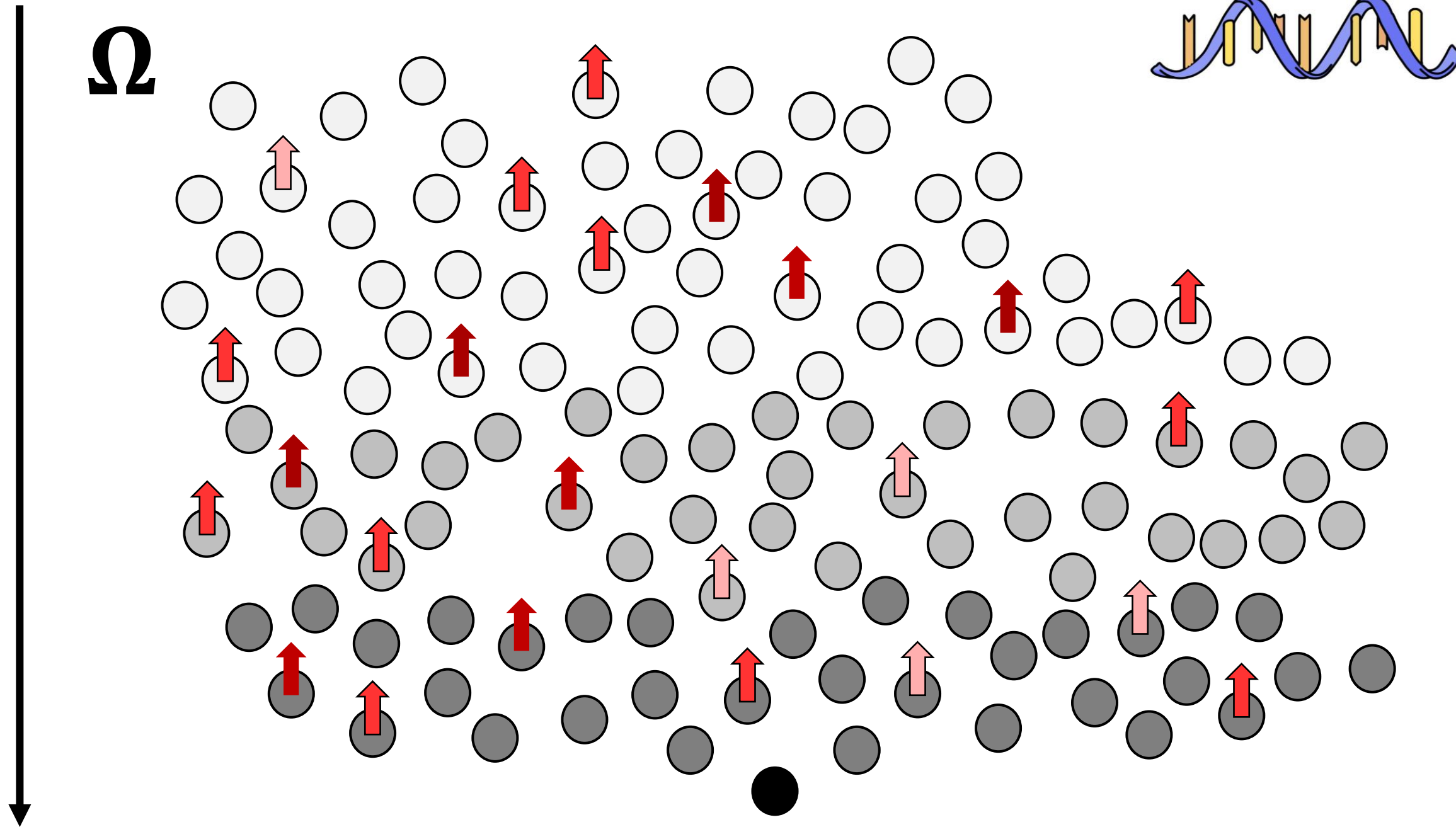
$\Delta G'(S)$



Ω

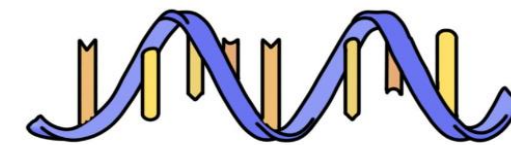
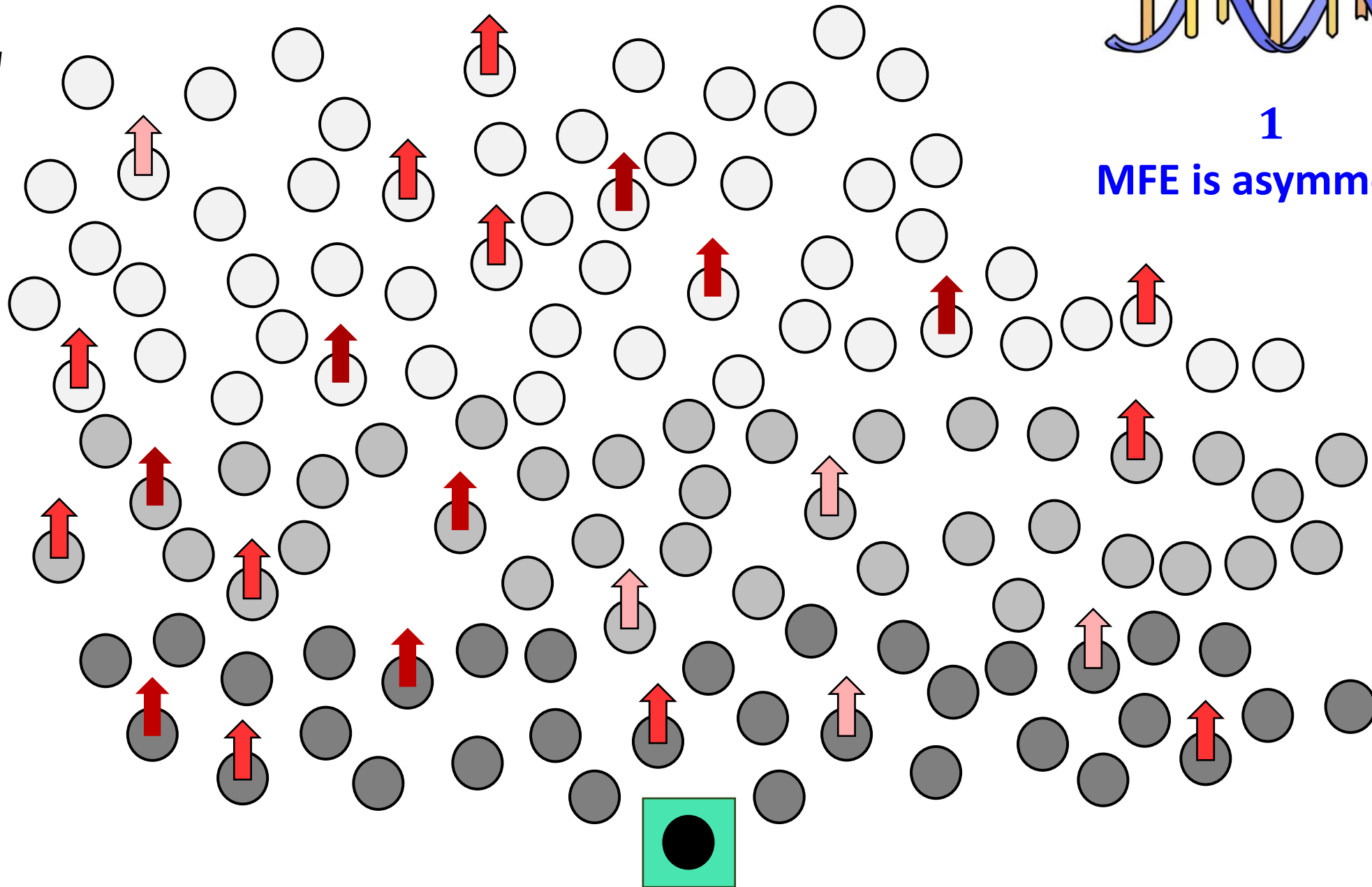
$\Delta G'(S)$



Ω $\Delta G'(S)$ 

Ω

$\Delta G'(S)$

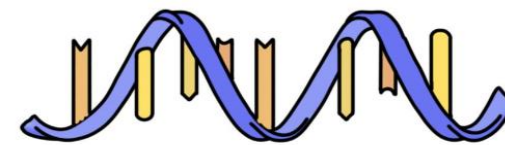
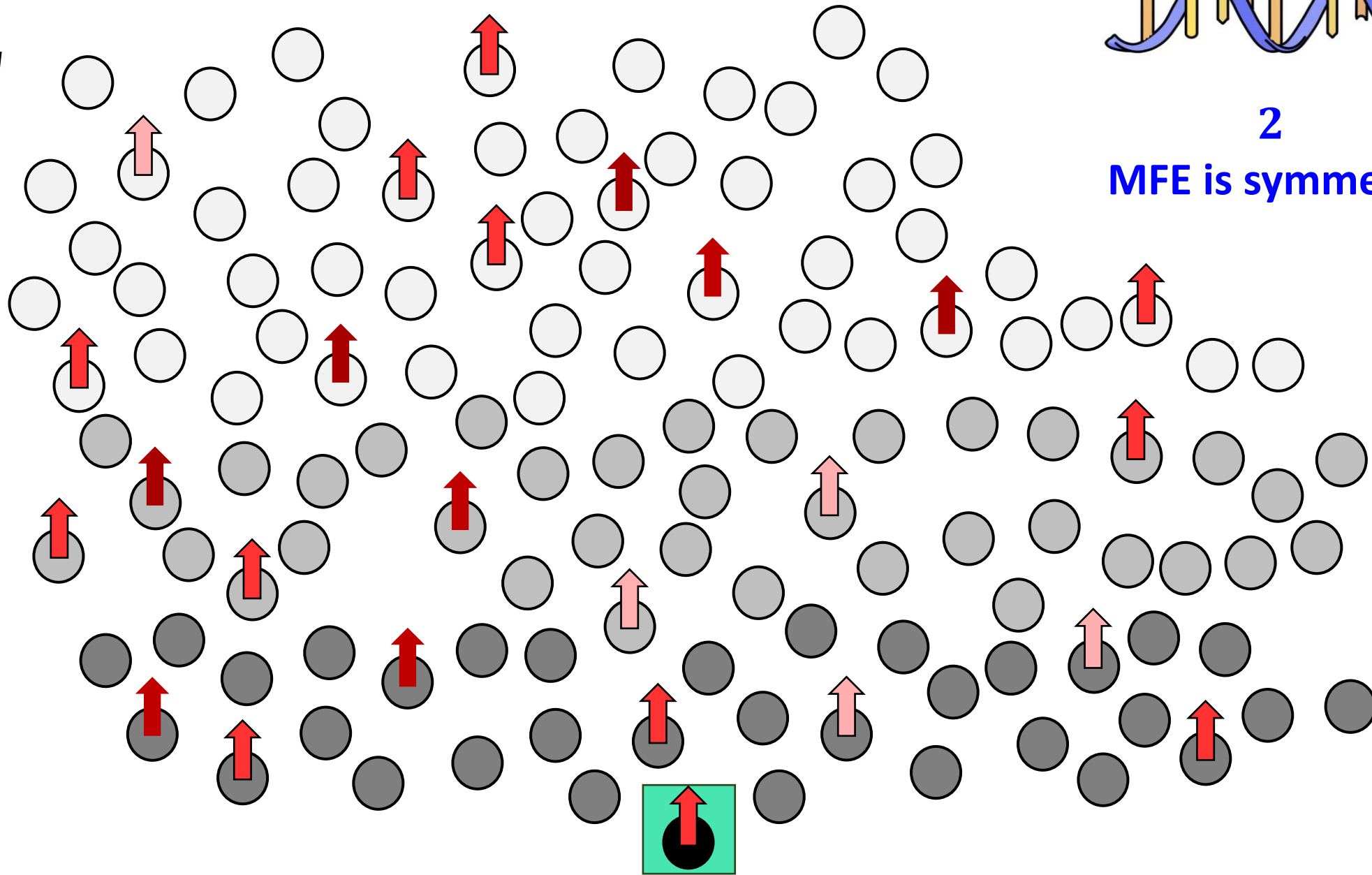


1

MFE is asymmetric

Ω

$\Delta G'(S)$



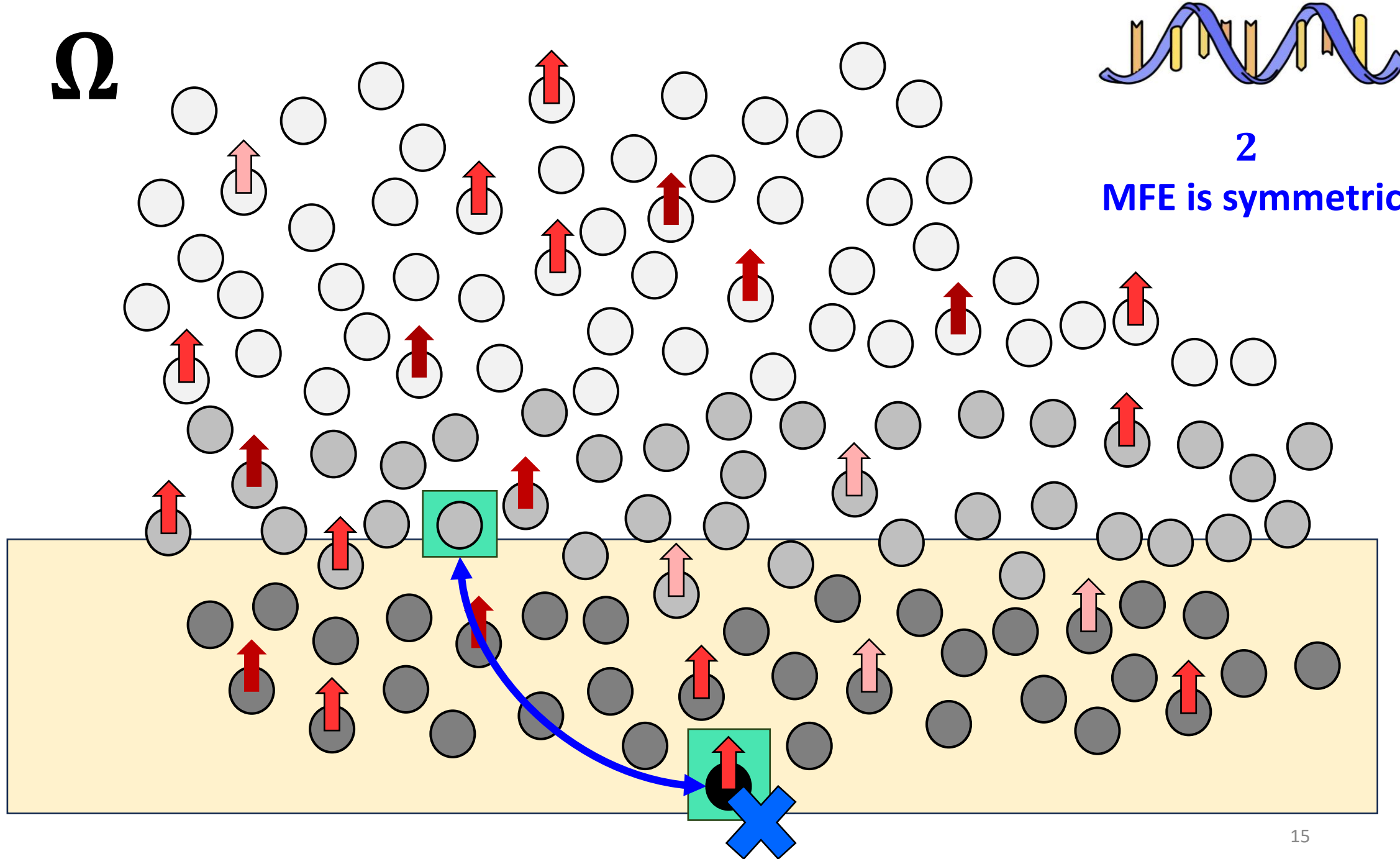
2

MFE is symmetric

$\Delta G'(S)$

$k_B T * \log R$

Ω

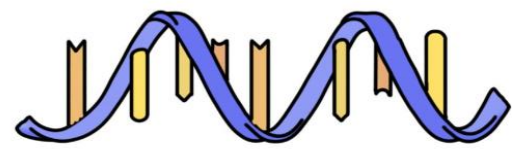
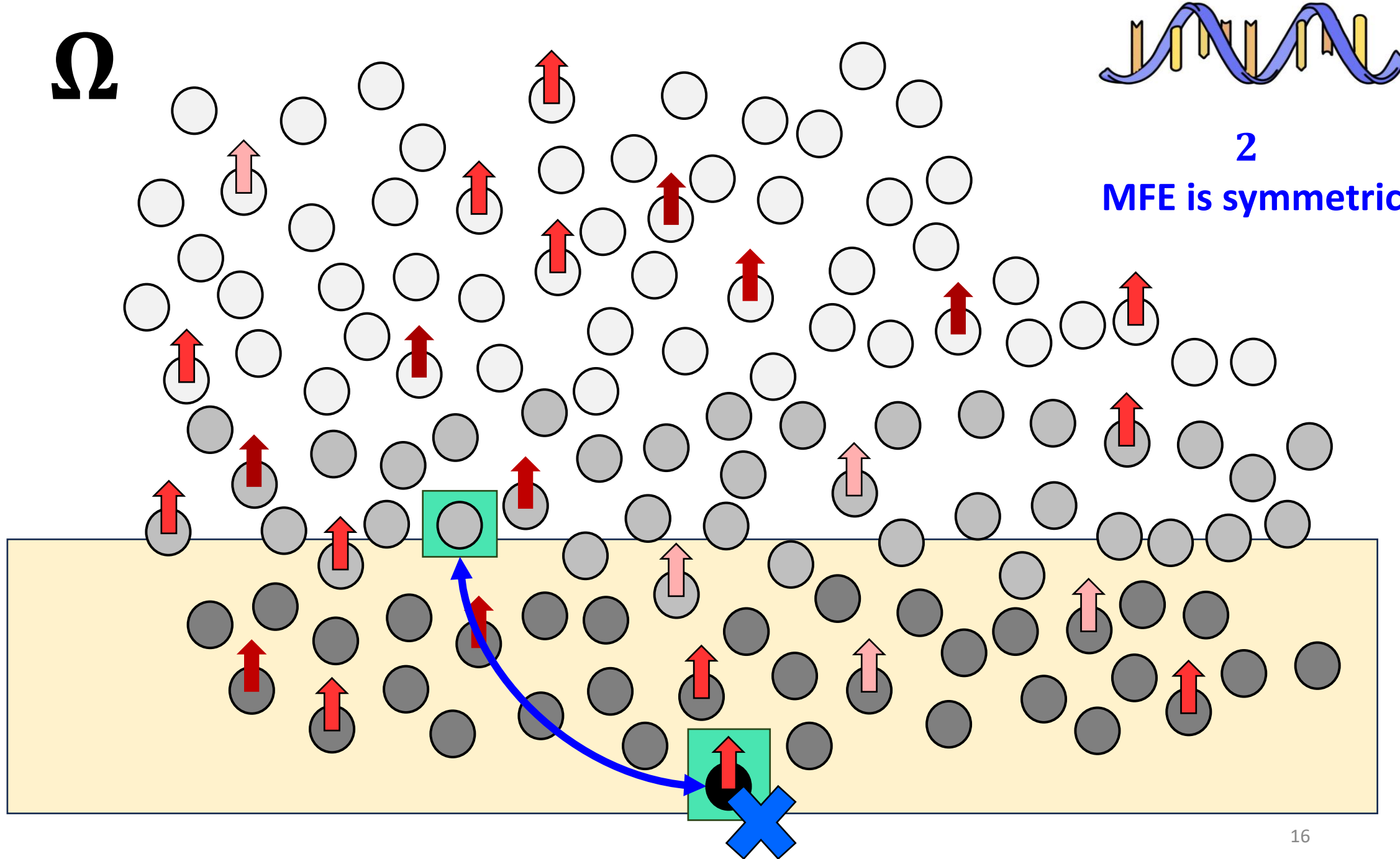


2
MFE is symmetric

$\Delta G'(S)$

Exponential

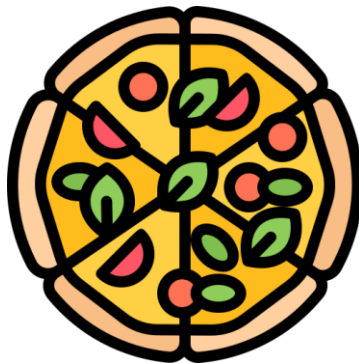
Ω

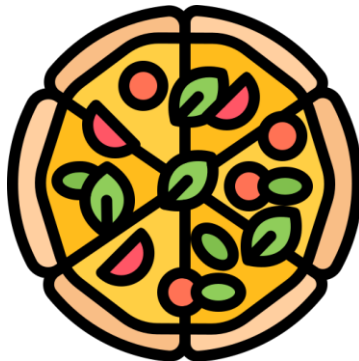


2

MFE is symmetric

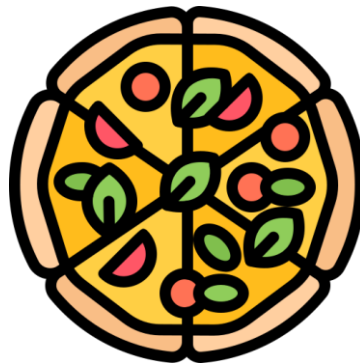
Our solution







S_x

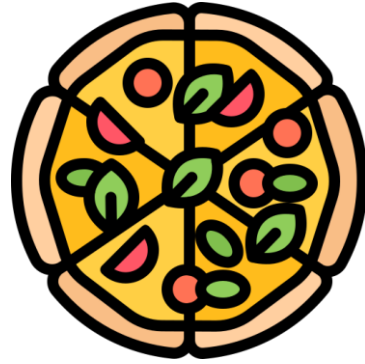


S_y

$\Delta G'(S)$



$$\Delta G'(S_y) \leq \Delta G'(S_x)$$



S_y

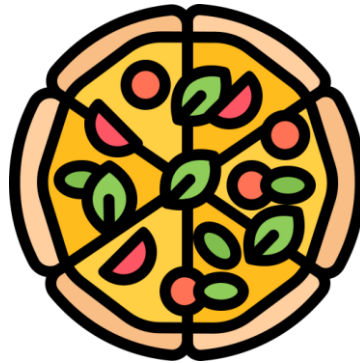


S_x

$\Delta G'(S)$



$$\Delta G'(S_y) \leq \Delta G'(S_x)$$



S_y Symmetric

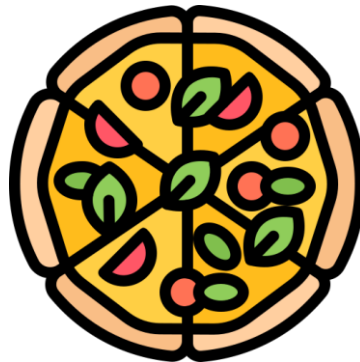


S_x Symmetric

$\Delta G'(S)$



$$\Delta G'(S_y) \leq \Delta G'(S_x)$$



S_y Symmetric



S_x Symmetric

S_x and S_y
Admissible cut

$\Delta G'(S)$



$$\Delta G'(S_y) \leq \Delta G'(S_x)$$

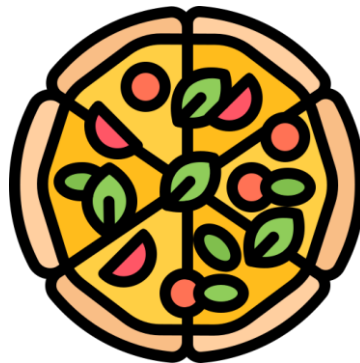


S_x Symmetric

X



S_z Asymmetric



S_y Symmetric

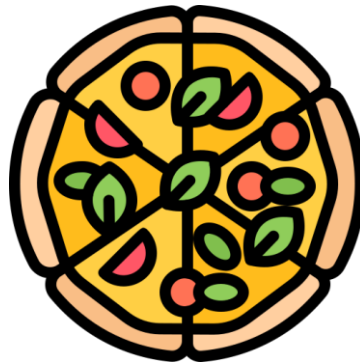
S_x and S_y
Admissible cut

$\Delta G'(S)$



$$\Delta G'(S_y) \leq \Delta G'(S_x)$$

$$\Delta G'(S_y) \leq G(S_z) \leq \Delta G'(S_x)$$



S_y Symmetric

X



S_z Asymmetric



S_x Symmetric

S_x and S_y
Admissible cut



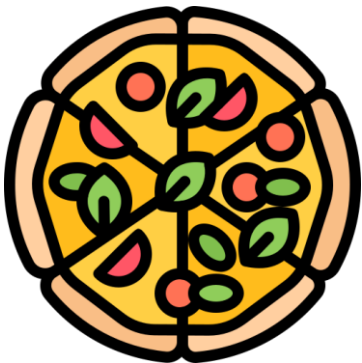
S_x
Symmetric

X

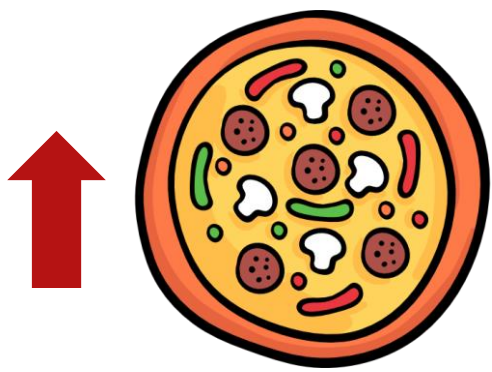


S_z
Asymmetric

S_x and S_y
Admissible cut



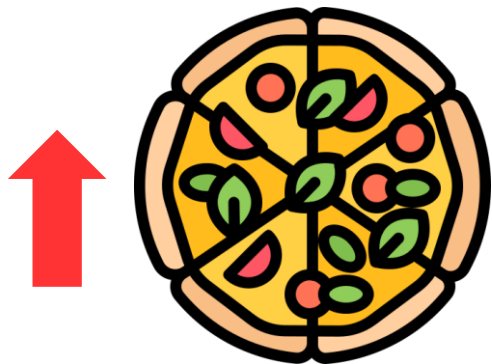
S_y
Symmetric



S_x
Symmetric

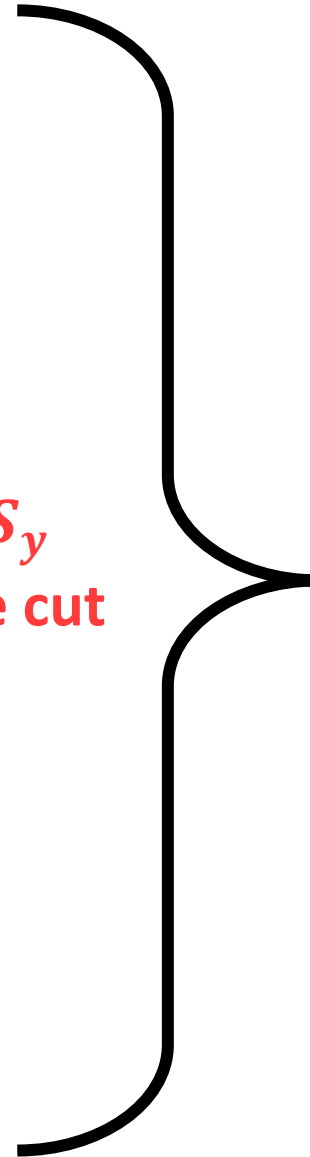


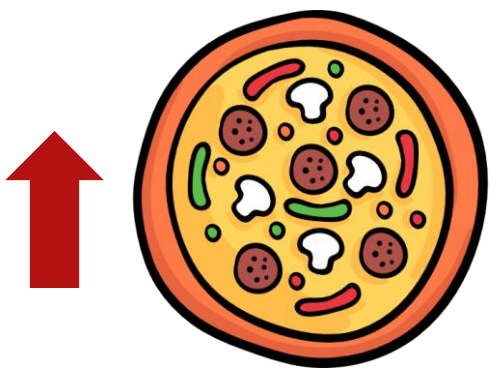
S_z
Asymmetric



S_y
Symmetric

S_x and S_y
Admissible cut



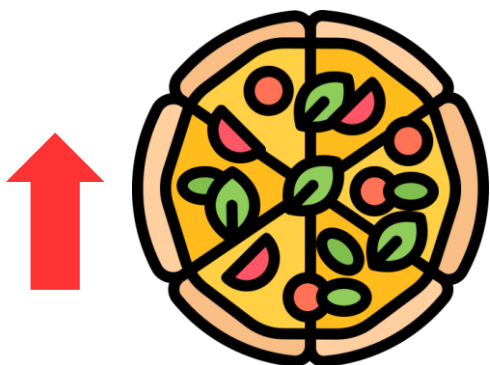


S_x
Symmetric



X

S_z
Asymmetric



S_y
Symmetric

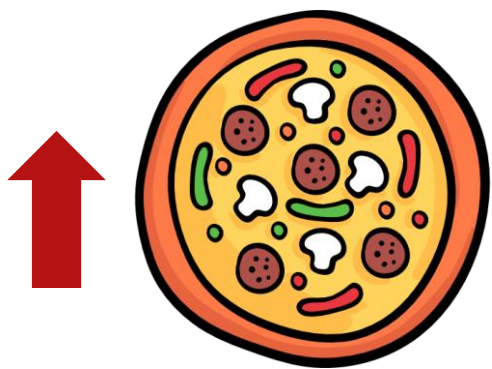
S_x and S_y
Admissible cut

Upper bound

$$\frac{N-c}{v(\pi)} (\sigma(v(\pi)) - v(\pi))$$

+

$$N^2/16$$

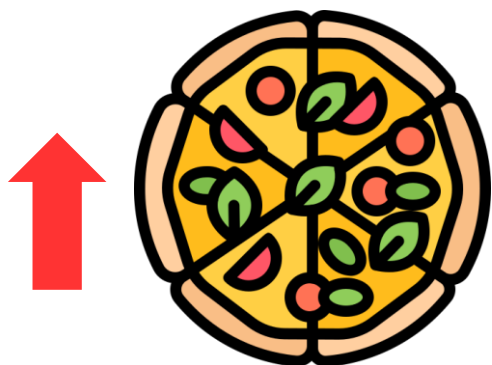


S_x
Symmetric



X

S_z
Asymmetric



S_y
Symmetric

S_x and S_y
Admissible cut

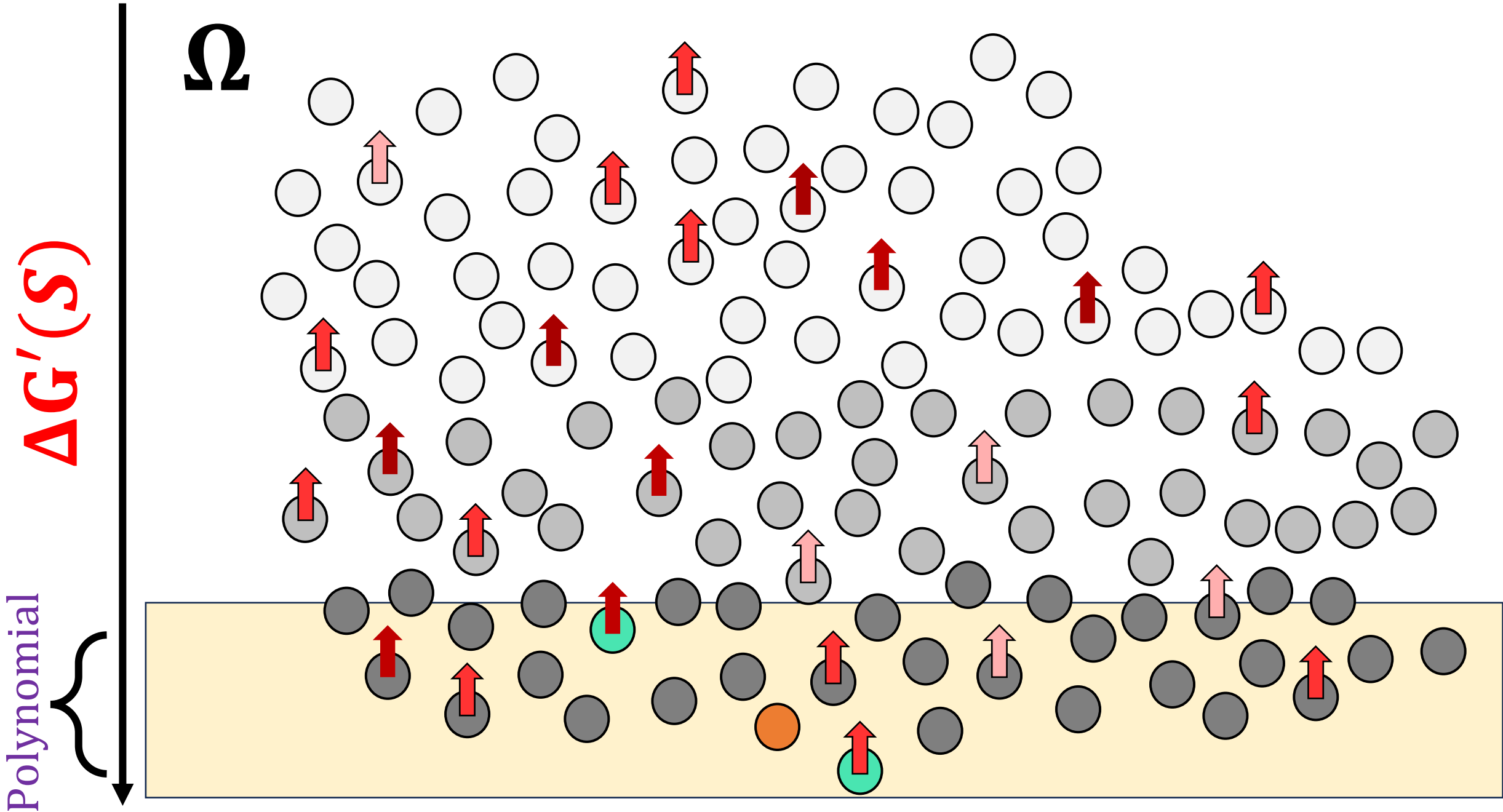
Upper bound

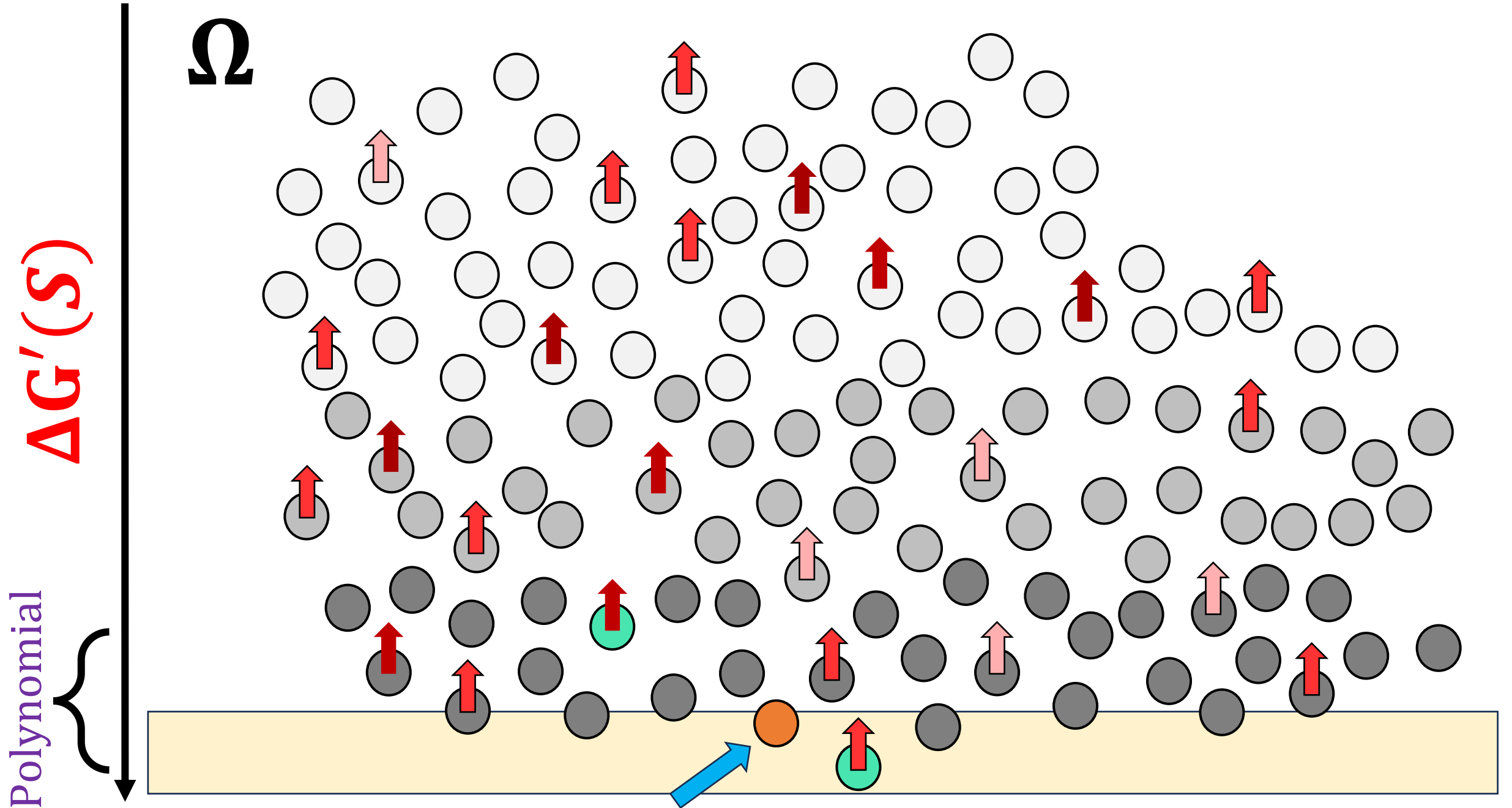
$$\frac{N-c}{v(\pi)} (\sigma(v(\pi)) - v(\pi))$$

+

$$N^2/16$$

Adjusting the backtracking algorithm to go through energy levels sequentially starting from the MFE level.





Computational complexity of Minimum Free Energy algorithms

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Thanks



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